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BINSEC/SE: A Dynamic Symbolic Execution Toolkit for Binary-level Analysis

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Abstract—When it comes to software analysis, several approaches exist from heuristic techniques to formal methods, which are helpful at solving different kinds of problems. Unfortunately very few initiative seek to aggregate these techniques in the same platform. BINSEC intend to fulfill this lack of binary analysis platform by allowing to perform modular analysis. This work focusses on BINSEC/SE, the new dynamic symbolic execution engine (DSE) implemented in BINSEC. We will highlight the novelties of the engine, especially in terms of interactions between concrete and symbolic execution or optimization of formula generation. Finally, two reverse engineering applications are shown in order to emphasize the tool effectiveness.

I. INTRODUCTION

While security can be checked and enforced at different level, we focus here on binary-level security analyses, such as malware comprehension or vulnerability analysis. Reverse engineering of binary codes is a key component of these activities, yet it is notoriously difficult. Especially, low-level assembly constructs can deceive disassemblers, while dynamic execution explores only a few possible behaviors. Our long-term goal is to adapt formal methods, which have been very successful in source-level safety analysis, to binary-level security analysis. We present in this paper a binary-level dynamic symbolic execution engine, named BINSEC/SE, geared toward security analysis and especially reverse-engineering.

The main contribution is a highly configurable generic DSE engine toolkit, with a strong interaction between the tracer and the symbolic execution core, as well as heavy optimizations on the path predicate. Two reverse engineering applications are shown to emphasize the tool effectiveness.

II. BACKGROUND

Dynamic Symbolic Execution [5], [9] is a formal technique for exploring program paths in a systematic way. For each path, the technique computes a path predicate, i.e. a set of constraints on the program input that leads to follow that path at runtime. This predicate is then fed to an automatic solver: a solution to the predicate is a new test input exploring the targeted path. Systematic exploration is achieved through iterating on all (user-bounded) paths of the program.

BINSEC [8] is a recent platform for formal analysis of binary codes. The platform currently proposes a front-end from \texttt{x86}(32bits) to a generic intermediate representation called DBA [8] (including decoding, disassembling, simplifications), a simulator and a static analysis engine (typically for dynamic jump resolution). The platform is written in OCaml (~30,000 loc), and will be released under an open source license soon.1 Our DSE engine BINSEC/SE is built upon BINSEC. The current tool paper describes only BINSEC/SE, a consequent add-on to BINSEC reusing only the x86-to-DBA translation.

III. USER VIEW

From a user point of view, the advantage of using a DSE engine is twofold. First, nowadays closed source binaries and malwares are so commonly widespread that having tools for working at binary level is mandatory. Second, recovering information from a binary can be a tedious task, especially in heavily obfuscated binaries. Therefore applying automatic techniques for recovering information is crucial with regard to the usual amount of code to review.

During a reverse-engineering task or an information retrieval task, the analyst can be interested in knowing information about the binaries, e.g. the register value at a given location, the possible values at a given memory address, whether a given branch can be covered or not or the possible targets of a dynamic jump. These problems can be addressed by DSE.

Various dynamic analyses are already implemented in BINSEC/SE and new ones can conveniently be written using a callback mechanism. An analysis takes a JSON configuration file as input and a trace file. Both the configuration and trace file formats are open and specified.

IV. BINSEC/SE ARCHITECTURE

We describe hereafter the new module for DSE, named BINSEC/SE. Its implementation is made of three components, presented in Figure 1.

- a pintool named PINSEC, written in C++ (~2,000 loc), based on the Pin DBI framework [10] for dynamic tracing;

1http://binsec.gforge.inria.fr
• the core DSE engine written in OCaml (∼8,000 loc), whose tasks are to generate path predicates, to send them to a solver and to get back new input data;
• the path selector, written in OCaml (∼1,200 loc), in charge of choosing which path to explore next.

B. DSE engine

Our DSE engine (Figure 1) follows a classical workflow. From an input trace, x86 instructions are translated into DBA instructions [8] (such a translation can be seen in Figure 2) on which the path predicate is computed [12], and then used to generate a formula (theory of arrays and bitvectors) which is then exported to the SMTLIB2 format\(^3\), the standard input of modern SMT solvers. A SMT solver is finally used to check the satisfiability of the formula and if so, to get back a solution (new input). We currently rely on Z3, CVC4 and Boolector. The connexion between BINSEC/SE and PINSEC easily allows to re-inject the new input to get a new trace.

Implementation insight. Internally an analysis is represented as an OCaml class, which each analysis should inherit from. This class provides appropriate callbacks and data structures for implementing any trace-based analysis. These callbacks allow to apply specific actions at a specific step and/or location along the execution in a highly configurable manner. They help developing new analyses without a deep understanding of the whole inner-working of the DSE. Main callbacks are:

- **pre_execution**, **post_execution** triggered respectively once at the beginning and the end of the analysis;
- **visit_instr_before**, **visit_instr_after** triggered before (resp. after) every asm instruction of the trace;
- **visit_dbainstr_before**, **visit_dbainstr_after** triggered before (resp. after) every DBA instructions of an x86 instruction;
- **input_message_received** triggered when a message is received from PINSEC.

A summary of all callback location calls is given in Figure 2.

C. Path selection

Beside a fine-grained control of a single trace execution, automatic multiple paths exploration is also a desired feature of a DSE tool. Ideally, the path exploration engine should allow either to fulfill some standard coverage requirements, or to focus on specific parts of the code through dedicated (user-defined) search heuristics. The design of the exploration module of BINSEC/SE is inspired from the one of OSMOSE [2]. It relies on a simple API allowing to easily implement various exploration strategies. In particular this API offers a function \(\text{select}(S)\) returning the “best” trace from a set of traces \(S\), w.r.t. a user-defined score function. Scores

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**A. The PINSEC tracer**

The tracer, called PINSEC, is a pintool [10] specialized in analyzing x86 architectures. Its main purpose is to generate the execution trace used as input by BINSEC. It is designed to offer a generic and modular tracing for both Linux and Windows binaries, and to export results into a generic format, here protobuf\(^2\) which is usable with major programming languages. The tracing is parameterized via a JSON file whose format is also defined in protobuf. The main parameters are:

- **start** and **stop**: addresses indicating where to begin and where to end the tracing,
- **call_skips**: address of calls for which the callee should not be traced,
- **fun_skips**: function addresses that should not be traced.

Other command line parameters allow to limit the trace size, the tracing time or the instrumentation scope.

**Advanced usage.** What is differentiating PINSEC from other pintools are the following functionalities:

- the automatic retrieval of function parameters and return values for (some) libraries functions;
- the possible retrieval of the concrete value of any register or memory location, even if not part of the operands / results of the currently instrumented instruction;
- the injection of symbolic or concrete values in any register or memory location, at any step of the execution;
- a remote command and control system allowing the DSE core and PINSEC to exchange messages in an interactive, debugger-like manner, in order to dynamically tune the instrumentation.

A full-fledged example using some of these parameters is presented in Section VI.

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\(^{2}\)https://developers.google.com/protocol-buffers/

\(^{3}\)http://smtlib.cs.uiowa.edu/
are built from several (predefined or dynamically computed) trace characteristics, such as length, last instruction call-depth, distance to a given target, etc. This approach allows to precisely define a wide-range of exploration strategies. Several strategies are already implemented such as DFS, BFS, random path, MinCall-DFS and MinCall-BFS [1].

V. FORMULA GENERATION AND SOLVING

The bottom-line of DSE is path predicate generation. While we follow the standard approach [12], i.e. maintaining a symbolic memory state along the computation together with the encountered path constraints (i.e. branching conditions), our method is original in several aspects:

• we can take into account constraints on the initial memory state (cf. Section V-B);
• we provide callback mechanisms for fine-tuning of concretization and symbolization (cf. Section V-C);
• the path predicate is highly optimized, and we take advantage of incremental solving (cf. Section V-A).

A. Optimizations

The core engine implements several optimizations of the path predicate, ranging from standard simplifications such as constant folding, local rewriting (e.g. \(a \oplus a = 0\)) or pruning useless parts of the formula, to less common optimizations, such as Read-over-Write (RoW) simplifications over array load and store [4]. Finally, great care is taken in order to make these optimizations compatible with the incremental solving mode of modern SMT solvers, which is particularly well-suited to DSE, since path predicates are naturally built incrementally.

Table I shows the results of optimization on a set of malware samples (cf. Section VI-B), while Table II shows a more detailed view on a single trace taken from the Artelad malware. The timeout is set to 20 seconds on a Intel Core i7 2.7 GHz (16GB RAM) using the Z3 solver.

<table>
<thead>
<tr>
<th>Malware</th>
<th>#Opt</th>
<th>#Operators</th>
<th>#VarsDef</th>
<th>Avg/query</th>
<th>Incremental</th>
</tr>
</thead>
<tbody>
<tr>
<td>Artelad</td>
<td>16K</td>
<td>1205</td>
<td>1290</td>
<td>28s/36</td>
<td>11m34</td>
</tr>
<tr>
<td>Benny</td>
<td>16K</td>
<td>5150</td>
<td>5110</td>
<td>18m37</td>
<td>7m34</td>
</tr>
<tr>
<td>Bogus</td>
<td>16K</td>
<td>191</td>
<td>191</td>
<td>41ms/1</td>
<td>145/298</td>
</tr>
<tr>
<td>Corned</td>
<td>944</td>
<td>52</td>
<td>51</td>
<td>1s/10</td>
<td>66ms/10</td>
</tr>
<tr>
<td>Esa</td>
<td>1600</td>
<td>115</td>
<td>113</td>
<td>3.44</td>
<td>1m59</td>
</tr>
<tr>
<td>Hiph</td>
<td>3252</td>
<td>1420</td>
<td>1410</td>
<td>10ms/23</td>
<td>11m34</td>
</tr>
<tr>
<td>Total</td>
<td>6193</td>
<td>1420</td>
<td>1410</td>
<td>11m34</td>
<td>11m34</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>time in seconds, or minutes/seconds</th>
<th>Avg/query: average time per query</th>
<th>Incremental: incremental solving turned off</th>
</tr>
</thead>
<tbody>
<tr>
<td>Optional optimizations</td>
<td>no-opt</td>
<td>pruning</td>
</tr>
<tr>
<td>#VarsDef</td>
<td>27373</td>
<td>8952</td>
</tr>
<tr>
<td>#Constraints</td>
<td>1585</td>
<td>1581</td>
</tr>
<tr>
<td>#Memory operations</td>
<td>1589</td>
<td>1585</td>
</tr>
<tr>
<td>#Operators</td>
<td>85805</td>
<td>6330</td>
</tr>
</tbody>
</table>

C. Concretization/Symbolization

For scalability issues, the path predicate is often approximated, either by forcing some logical values to their runtime values (observed by the tracer) – concretization, or by injecting fresh logical input – symbolization. Both operations reduce the complexity of the formula, at the price of completeness and/or correctness. So choosing appropriately when and what to concretize/symbolize is a major practical issue. While most tools offer only hardcoded policies, BINSEC provides a complete set of callbacks to specify such choices. Such a mechanism allows, for instance, symbolic reasoning on memory pointers, where existing tools usually enforce memory addresses to be concrete.

VI. EXPERIMENTATIONS

A. Reverse engineering

In reverse-engineering, so-called crackme challenges simulate real world situations where binary programs are designed to make the analysis difficult. The goal is usually to find a string key (the flag) allowing to validate the challenge.

Flare-On is a reverse engineering challenge organized since 2014 by FireEye Security⁴. We analyse the first challenge of 2015 since it is straightforward enough to be discussed in detail. This crackme is a 32 bits Windows program which asks for a password and prints “You are success” or “You are failure” depending on the keyboard input. Each byte of the input (input_buffer) is xorred with 0x7d, the result is then checked against a key stored in the data section (data_str), cf. Figure 3.

![Fig. 3. Flare-on #1 key loop computation](http://www.flare-on.com/)

Solving the crackme by symbolic execution aims at checking that the predicate 2F=1 is true at location 0x40105b; and this, for every character of the key. If the generated formula is
satisfiable, then the right character can simply be retrieved
in the formula solution at \texttt{input\_buffer[ecx]}. One last
remaining problem is the initial state. If none is specified,
the solver is allowed to give any valuation to the content
of \texttt{data\_str}, while it is not user-controllable. Hence, an
initial state constraining the value of the \texttt{data\_str} bytes is
required in order to get meaningful results for \texttt{input\_buffer}
(cf. Section V-B). Finally, iterating over \texttt{input\_buffer} can be
done in two ways:

1) From a trace taking the fail branch, we should solve
\(ZF=1\) at 0x40105b to get the right char and inject it
back as input via the configuration file for generating a
new trace that will do a second loop iteration, repeating
until the whole key is found;

2) Configure a breakpoint at 0x40105b (the address of
\texttt{jnz}), compute the right character value, send a com-
mand to patch the ZF flag and resume to force the
tracer to take the right branch and looping again. This
method is fully automatic and allow to take advantage
of incremental solving (cf. Section V-A).

We obtain the right key solving the challenge,
bunny_sl0pe@flare-on.com\(^5\). Note that we have solved
several other Flare-on challenges with the same principle.

B. Malware exploration

We are interested here in demonstrating the ability of
BINSEC/SE to automatically discover new code areas of a
program, which is especially crucial in malware analysis. We
consider 11 malware programs taken from the VX Heaven
database\(^6\) and used in recent works on deobfuscation [11].
Table III shows our results. We report the number of new
explored paths (i.e. each time DSE manages to negate a branch
of the original trace), and of new behaviors (i.e. each time DSE
covers a branch not yet covered by the initial trace or by a
generated one). Here, we are able to discover 43 new code
areas in the malwares under consideration.

### Table III. Malware Exploration

<table>
<thead>
<tr>
<th>Path</th>
<th>Brth</th>
<th>Mt</th>
<th>Br</th>
<th>Fmt</th>
<th>Avg</th>
<th>Total</th>
</tr>
</thead>
<tbody>
<tr>
<td>Virus.Win32.Antelad.2173</td>
<td>10k</td>
<td>1295</td>
<td>3</td>
<td>652</td>
<td>2</td>
<td>0.59</td>
</tr>
<tr>
<td>Virus.Win32.Bellad.2</td>
<td>10k</td>
<td>1551</td>
<td>3</td>
<td>602</td>
<td>2</td>
<td>0.30</td>
</tr>
<tr>
<td>Virus.Win32.Benny.3194.a</td>
<td>10k</td>
<td>3153</td>
<td>121</td>
<td>4</td>
<td>4</td>
<td>0.6</td>
</tr>
<tr>
<td>Virus.Win32.Bojag.4096</td>
<td>10k</td>
<td>191</td>
<td>1</td>
<td>0</td>
<td>0</td>
<td>0.02</td>
</tr>
<tr>
<td>Virus.Win32.Callistino.4</td>
<td>10k</td>
<td>2049</td>
<td>819</td>
<td>1</td>
<td>1</td>
<td>0.08</td>
</tr>
<tr>
<td>Virus.Win32.Comat.2</td>
<td>804</td>
<td>12</td>
<td>17</td>
<td>12</td>
<td>12</td>
<td>0.63</td>
</tr>
<tr>
<td>Virus.Win32.Feva.a</td>
<td>1099</td>
<td>115</td>
<td>49</td>
<td>4</td>
<td>4</td>
<td>0.04</td>
</tr>
<tr>
<td>Virus.Win32.Horpa</td>
<td>2152</td>
<td>1105</td>
<td>9</td>
<td>508</td>
<td>6</td>
<td>0.49</td>
</tr>
<tr>
<td>Virus.Win32.Paljirr.6</td>
<td>10k</td>
<td>1520</td>
<td>9</td>
<td>508</td>
<td>6</td>
<td>2.23</td>
</tr>
<tr>
<td>Virus.Win32.Seppusa.1606</td>
<td>1696</td>
<td>147</td>
<td>61</td>
<td>4</td>
<td>4</td>
<td>0.05</td>
</tr>
<tr>
<td>Virus.Win32.Wd.a</td>
<td>10k</td>
<td>3358</td>
<td>2</td>
<td>3358</td>
<td>2</td>
<td>0.00</td>
</tr>
<tr>
<td>Total</td>
<td>7360</td>
<td>1019</td>
<td>764</td>
<td>597</td>
<td>45</td>
<td>8.17</td>
</tr>
</tbody>
</table>

\texttt{br}: number of branches – \texttt{uniq}: unique conditional jumps – \texttt{path}: number of
new paths – \texttt{behavior}: new path covering a new code area – \texttt{Avg}: Average
solving time per path

VII. RELATED WORK

A few other initiatives aim at applying DSE on binary pro-
grams for security purposes, most notably BAP [3], S2E [7],
FuzzBall [6] or Triton [13]. All these approaches also rely
on an intermediate representations similar to DBA. Regarding
dynamic tracing, these tools are built either on Qemu or on Pin.
In addition to a generic DSE engine, the novelties provided
by BINSEC/SE are highly configurable path exploration strategies
and concretization/symbolization policies, a strong interaction
between the tracer and the DSE core and an optimized path
predicate computation.

VIII. CONCLUSION

The dynamic symbolic execution implemented in the plat-
form yields the interesting property of being modular w.r.t
the analysis configuration and to provide a strong interaction
between the DSE (BINSEC) and the tracer (PINSEC). Some
functionalties like the formula optimizations makes it original
from the existing state-of-the-art tools. This platform comes
really handy for reverse-engineering and information recovery
tasks like it is the case for source-less programs and malwares.

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\(^5\) A complete demo is available at https://youtu.be/0xUc2jbpojQo
\(^6\)http://vxheaven.org/